

# Identifiability in Sparse Factor Analysis

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## Spearman's One-Factor Model (1904)



Linear structural equations:

$$X_{1} = \lambda_{10} + \lambda_{1L}L + \varepsilon_{1},$$
  

$$X_{2} = \lambda_{20} + \lambda_{2L}L + \varepsilon_{2},$$
  

$$X_{3} = \lambda_{30} + \lambda_{3L}L + \varepsilon_{3},$$
  

$$X_{4} = \lambda_{40} + \lambda_{4L}L + \varepsilon_{4}.$$

Jointly independent errors:  $\varepsilon_1, \ldots, \varepsilon_4$ .

$$\operatorname{Var}[\varepsilon_{v}] = \omega_{vv} < \infty$$
,  $\operatorname{Var}[L] = 1$ .

Topic: Can we recover the "factor loadings" 
$$\lambda_{vL}$$
 and the "error variances"  $\omega_{vv}$  from  $\Sigma = Var[X]$ ?

### Socio-economic Example from Harmann (1976), Modern Factor Analysis





$$X = \Lambda L + \varepsilon, \quad \text{where } \Lambda = \begin{pmatrix} \lambda_{11} & 0 \\ 0 & \lambda_{22} \\ \lambda_{31} & 0 \\ \lambda_{41} & \lambda_{42} \\ \lambda_{51} & \lambda_{52} \end{pmatrix}$$

Observed covariance matrix (when Var[L] = I):

$$\boldsymbol{\Sigma} = \boldsymbol{\Lambda}\boldsymbol{\Lambda}^{\top} + \boldsymbol{\Omega} = \begin{pmatrix} \omega_{11} + \lambda_{11}^2 & 0 & \lfloor \lambda_{11}\lambda_{31} \rfloor & \lfloor \lambda_{11}\lambda_{41} \rfloor & \lambda_{11}\lambda_{51} \\ 0 & \omega_{22} + \lambda_{22}^2 & 0 & \lambda_{22}\lambda_{42} & \lambda_{22}\lambda_{52} \\ \lambda_{11}\lambda_{31} & 0 & \omega_{33} + \lambda_{31}^2 & \lfloor \lambda_{31}\lambda_{41} \rfloor & \lambda_{31}\lambda_{51} \\ \lambda_{11}\lambda_{41} & \lambda_{22}\lambda_{42} & \lambda_{31}\lambda_{41} & \omega_{44} + \lambda_{41}^2 + \lambda_{52}^2 & \lambda_{41}\lambda_{51} + \lambda_{42}\lambda_{52} \\ \lambda_{11}\lambda_{51} & \lambda_{22}\lambda_{52} & \lambda_{31}\lambda_{51} & \lambda_{41}\lambda_{51} + \lambda_{42}\lambda_{52} & \omega_{55} + \lambda_{51}^2 + \lambda_{52}^2 \end{pmatrix}$$

We see that

1)  $\sqrt{\frac{\sigma_{13}\sigma_{14}}{\sigma_{34}}} = \sqrt{\frac{\lambda_{11}\lambda_{31}\lambda_{11}\lambda_{41}}{\lambda_{31}\lambda_{41}}} = \sqrt{\lambda_{11}^2} = a_1\lambda_{11}$  with  $a_1 \in \{\pm 1\}$  and  $\sigma_{34} = \lambda_{31}\lambda_{41} \neq 0$  'almost surely', 2)  $\frac{\sigma_{13}}{\sqrt{\sigma_{13}\sigma_{14}/\sigma_{34}}} = \frac{\lambda_{11}\lambda_{31}}{a_1\lambda_{11}} = a_1\lambda_{31}$  with  $\lambda_{11} \neq 0$  'almost surely'.  $\implies$  Can identify  $\Lambda_{ch(L_1),L_1}$  up to column-sign, similarly  $\Lambda_{ch(L_2),L_2}$ .

## Factor Analysis vs. Causal Representation Learning



Understanding sparse factor analysis is key for causal representation learning!





#### Variables:

Observed: 
$$X = (X_v)_{v \in V}$$
 Latent:  $L = (L_h)_{h \in H}$ 

Graph:

Bipartite directed graph 
$$G = (V \dot{\cup} \mathcal{H}, D)$$
, that is,  $D \subseteq \mathcal{H} \times V$ .

Sparse factor analysis model:

 $X = \Lambda L + \varepsilon$ 

- all latent factors and error terms in  $(L, \varepsilon)$  are mutually independent, so  $\Omega = Var[\varepsilon] = diag(\omega_{vv} : v \in V)$  diagonal, and Var[L] = I.
- parameter matrix  $\Lambda$  is sparse and supported over edge set D (write  $\Lambda \in \mathbb{R}^D$ ).

### Content of the Talk

### Definition

Every factor analysis graph G yields a parametrization of the observed covariance matrix:

$$\tau_{G} : (\Lambda, \Omega) \longmapsto \Sigma \equiv \Lambda \Lambda^{\top} + \Omega$$
  
Fiber:  $\mathcal{F}_{G}(\Omega, \Lambda) = \{ (\widetilde{\Omega}, \widetilde{\Lambda}) : \tau_{G}(\widetilde{\Omega}, \widetilde{\Lambda}) = \tau_{G}(\Omega, \Lambda) \}.$ 

The model given by G is generically sign-identifiable if

$$\mathcal{F}_{G}(\Omega,\Lambda) = \{ (\widetilde{\Omega},\widetilde{\Lambda}) : \widetilde{\Omega} = \Omega \text{ and } \widetilde{\Lambda} = \Lambda \Psi \text{ for } \Psi \in \{\pm 1\}^{|\mathcal{H}| \times |\mathcal{H}|} \text{ diagonal} \} \qquad \text{for `almost all'} (\Lambda,\Omega).$$

#### Main Contribution:

- Sufficient graphical condition for generic sign-identifiability.
- Recursive polynomial time algorithm.

(caveat: polynomial time when bounding a cardinality in a search step)

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Gröbner basis computations solve the problem ... on a small scale.



## Our Software



```
# Define graph
> L = matrix(c(1, 0,
                                                                       (Employ)
                                                               (School)
                                                                                (Service)
                                                                                        (House)
                                                         Por
                0, 1,
+
                1, 0,
+
                1, 1,
+
                1, 1), 5, 2, byrow=TRUE)
+
> g = FactorGraph(L)
>
> # Check identifiability
> res = mID(g)
Generic Sign-Identifiability Summary
# nr. of latent nodes that are gen. sign-identifiable: 2
# gen. sign-identifiable nodes: 1, 2
```

Available at https://github.com/MiriamKranzlmueller/id-factor-analysis.

### Rotational Indeterminacy



• No restriction on  $\Lambda$ :

$$\Lambda\Lambda^{\top} + \Omega = (\Lambda Q)(Q^{\top}\Lambda^{\top}) + \Omega$$
 for all  $Q$  orthogonal.

- Literature: Focus on identifiability of  $\Omega$  and  $\Lambda\Lambda^{\top}$  in full factor analysis models.
- However, if  $\Lambda$  is sparse, we might get generic sign-identifiability.

## Zero Upper Triangular Assumption



### Definition

The graph G satisfies the Zero Upper Triangular Assumption (ZUTA) if there exists an ordering  $\prec$  on the latent nodes  $\mathcal{H}$  such that ch(h) is not contained in  $\bigcup_{\ell \succ h} ch(\ell)$  for all  $h \in \mathcal{H}$ .

### Example

	$h_1$	$h_2$	h <sub>3</sub>			$h_1$	$h_3$	$h_2$			$h_1$	$h_3$	$h_2$
$v_1$	(*	0	0)		$v_1$	(*	0	0)		$v_1$	(*	0	0)
<b>v</b> <sub>2</sub>	*	*	*	ZUTA	<b>v</b> <sub>2</sub>	*	*	*	upper-tri=0 ~→ upper-diag≠0	V <sub>3</sub>	*	*	0
<i>V</i> 3	*	0	*		V <sub>3</sub>	*	*	0		<b>v</b> <sub>2</sub>	*	*	*
<i>V</i> 4	*	*	*		<i>V</i> 4	*	*	*		<i>V</i> 4	*	*	*
<i>V</i> 5	*	*	*		$V_5$	*	*	*		<b>V</b> 5	*	*	*
V <sub>6</sub>	0/	*	*/		V <sub>6</sub>	0/	*	*/		<i>V</i> 6	0/	*	*/

ZUTA = "can permute cols and rows such that the upper right triangle of  $\Lambda$  is zero" and w.l.o.g. "diagonal entries are nonzero".

 $V_5$ 

 $V_6$ 

### Anderson-Rubin Criterion



Theorem [Anderson, Rubin (1956)]

A factor analysis graph  $G = (V \cup \mathcal{H}, D)$  that satisfies ZUTA is generically sign-identifiable if for any deleted row of the symbolic matrix  $\Lambda = (\lambda_{vh}) \in \mathbb{R}^D$  there exist two disjoint submatrices that are generically of rank  $|\mathcal{H}|$ .

### Example



$$\Lambda = \begin{pmatrix} \lambda_{v_1h_1} & 0\\ \lambda_{v_2h_1} & \lambda_{v_2h_2}\\ \lambda_{v_3h_1} & \lambda_{v_3h_2}\\ 0 & \lambda_{v_4h_2}\\ 0 & \lambda_{v_5h_2} \end{pmatrix}$$

**Observation**:

Need  $|V| \geq 2|\mathcal{H}| + 1$ .

### Examples for Inconclusiveness of Anderson-Rubin



1)



[Hosszejni, Frühwirth-Schnatter (2022)]

2)

 $\lambda_{v_1h_1}$ 0  $\lambda_{v_2 h_1}$  $\lambda_{v_2h_2}$  $\lambda_{v_3h_1}$  $\lambda_{v_3h_2} \lambda_{v_3h_3}$  $\lambda_{v_4 h_1}$  $\lambda_{v_4h_2}$  $\lambda_{v_4h_3}$  $\lambda_{v_5 h_1}$  $\lambda_{v_5 h_3}$  $\lambda_{V_5h_2}$  $\lambda_{v_6h_1}$ 0  $\lambda_{v_6h_2}$ 

3)

0  $\lambda_{v_1h_1}$  $\lambda_{v_2h_1}$ 0 0  $\lambda_{v_2h_2}$ 0  $\lambda_{v_3h_2}$ 0  $\lambda_{v_3h_3}$  $\lambda_{v_4h_1}$ 0  $\lambda_{v_4 h_3}$  $\lambda_{v_4h_4}$ 0 0 0  $\lambda_{v_5 h_2}$ 0 0 0 0 0 0 0  $\lambda_{v_6 h_3}$ 0  $\lambda_{v_7 h_4}$ 0 0 0  $\lambda_{v_8h_4}$ 0  $\lambda_{v_9 h_4/2}$ 

### Investigation of Anderson-Rubin



1) Anderson-Rubin: Fix  $v \in V$ . Find disjoint  $U, W \subseteq V \setminus \{v\}$  with  $|U| = |W| = |\mathcal{H}|$  such that  $\det(\Lambda_{U,\mathcal{H}}) \neq 0$  and  $\det(\Lambda_{W,\mathcal{H}}) \neq 0$  (not the zero polynomial)  $\iff \det([\Lambda\Lambda^{\top}]_{U,W}) \neq 0.$ 

2) Consider the matrix with exactly one diagonal entry:

$$[\Lambda\Lambda^{\top}]_{\{\nu\}\cup U,\{\nu\}\cup W} = \left(\frac{[\Lambda\Lambda^{\top}]_{\nu\nu} \mid [\Lambda\Lambda^{\top}]_{\nu,W}}{[\Lambda\Lambda^{\top}]_{U,\nu} \mid [\Lambda\Lambda^{\top}]_{W,U}}\right) = \left(\frac{[\Lambda\Lambda^{\top}]_{\nu\nu} \mid \Sigma_{\nu,W}}{\Sigma_{U,\nu} \mid \Sigma_{U,W}}\right).$$

3) Solve for diagonal entry  $[\Lambda\Lambda^{\top}]_{\nu\nu}$  by Laplace expansion:

$$0 = \det([\Lambda\Lambda^{\top}]_{\{v\}\cup U, \{v\}\cup W}) = [\Lambda\Lambda^{\top}]_{vv} \underbrace{\det(\Sigma_{U,W})}_{\neq 0} - \sum_{w \in W} \operatorname{sign}(w)\sigma_{vw} \det(\Sigma_{U,\{v\}\cup W \setminus \{w\}}).$$

4) Conclude: Solving for diagonal entries of  $\Lambda\Lambda^{\top}$  is equivalent to solving for  $\Omega$ .

 $\implies$  Equivalent to solving for  $\Sigma - \Omega = \Lambda \Lambda^{\top}$ .

 $\stackrel{\text{ZUTA}}{\Longrightarrow}$  generic sign-identifiability (uniqueness of Cholesky decomposition).

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### Main Idea for Sparse Setup

Local approach: Can also choose U, W such that  $|W| = |U| < |\mathcal{H}|$ .  $\rightsquigarrow$  Ensure that det $([\Lambda\Lambda^{\top}]_{U,W}) \neq 0$  and that det $([\Lambda\Lambda^{\top}]_{\{v\}\cup U,\{v\}\cup W}) = 0$ .

**Characterization**: When is det( $[\Lambda \Lambda^{\top}]_{A,B}$ ) = 0 if  $\Lambda$  is sparse?

 $\rightsquigarrow$  Intersection-free matchings.

#### Definition

System of paths  $\Pi = \{\pi_1, \dots, \pi_k\}$  is matching of  $A = \{a_1, \dots, a_k\} \subseteq V$  and  $B = \{b_1, \dots, b_k\} \subseteq V$  if  $\pi_i = a_i \leftarrow h_i \rightarrow b_i$ .

A matching is intersection-free if all latent nodes  $h_i$  are distinct.

#### Lemma

For two subsets  $A, B \subseteq V$  with |A| = |B| it holds that  $det([\Lambda\Lambda^{\top}]_{A,B}) \neq 0$  if and only if there is an intersection-free matching of A and B.

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## Matching Criterion



#### Definition

Fix a latent node  $h \in \mathcal{H}$ . Tuple  $(v, U, W, S) \in V \times 2^V \times 2^V \times 2^{\mathcal{H} \setminus \{h\}}$  satisfies the matching criterion for h if

(i)  $pa(v) \setminus S = \{h\}$  and  $v \notin U \cup W$ ,

(ii) U and W are disjoint, nonempty sets of equal cardinality,

(iii) there exists an intersection-free matching of U and W that avoids S,

(iv) there does not exist an intersection-free matching of  $\{v\} \cup W$  and  $\{v\} \cup U$  that avoids S.

S = "solved nodes".

By (iii), det( $[\Lambda\Lambda^{\top}]_{U,W}$ )  $\neq 0$ . By (iv), det( $[\Lambda\Lambda^{\top}]_{\{v\}\cup U,\{v\}\cup W}$ ) = 0.

# Algorithm: Recursive Solving



### Theorem (M-identifiability)

If the tuple (v, U, W, S) satisfies the matching criterion with respect to h and all nodes  $\ell \in S$  are "solved before", then we can "solve" for h.

That is,  $(\widetilde{\Omega}, \widetilde{\Lambda}) \in \mathcal{F}_{G}(\Omega, \Lambda) \implies \widetilde{\Lambda}_{ch(h),h} = \pm \Lambda_{ch(h),h}.$ 

### Algorithm (M-ID)

- Cycle through latent nodes h and search for tuples (v, U, W, S).
- Network-flow setup finds suitable tuples in polynomial time under a bound on |U| = |W|.

Conjecture: If we do not bound the cardinality |U| = |W|, then M-ID is NP-complete.

### Remarks

- Subsumes Anderson-Rubin.
- Together with an extension, subsumes anything we know (e.g. Bekker and ten Berge, 1997).

Example



$$\Lambda = \begin{pmatrix} \lambda_{v_1h_1} & 0 & 0 \\ \lambda_{v_2h_1} & \lambda_{v_2h_2} & 0 \\ \lambda_{v_3h_1} & \lambda_{v_3h_2} & \lambda_{v_3h_3} \\ \lambda_{v_4h_1} & \lambda_{v_4h_2} & \lambda_{v_4h_3} \\ \lambda_{v_5h_1} & \lambda_{v_5h_2} & \lambda_{v_5h_3} \\ \lambda_{v_6h_1} & \lambda_{v_6h_2} & 0 \end{pmatrix}$$

$$\begin{array}{ll} h_1: \text{ Take } v = v_1, \ S = \emptyset \text{ and } U = \{v_2, v_6\}, \ W = \{v_3, v_4\}.\\ (\text{iii)} \ v_2 \leftarrow h_1 \rightarrow v_3, \ v_6 \leftarrow h_2 \rightarrow v_4; \quad (\text{iv)} \ \mathsf{pa}(\{v\} \cup U) \cap \mathsf{pa}(\{v\} \cup W) = \{h_1, h_2\}. \end{array}$$

 $\begin{array}{l} h_2: \mbox{ Take } v = v_2, \ S = \{h_1\} \ \mbox{and } U = \{v_3\}, \ W = \{v_6\}. \\ (\mbox{iii) } v_3 \leftarrow h_2 \rightarrow v_6; \quad (\mbox{iv) } (pa(\{v\} \cup U) \cap pa(\{v\} \cup W)) \setminus S = \{h_2\}. \end{array}$ 

$$\begin{array}{l} h_3: \text{ Take } v = v_3, \ S = \{h_1, h_2\} \text{ and } U = \{v_4\}, \ W = \{v_5\}.\\ (\text{iii}) \ v_4 \leftarrow h_3 \rightarrow v_5; \quad (\text{iv}) \ (\mathsf{pa}(\{v\} \cup U) \cap \mathsf{pa}(\{v\} \cup W)) \setminus S = \{h_3\}. \end{array}$$

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### Conclusion

- Sparse factor analysis models are the "building block" for many latent variable models.
- Latent variable models generally feature complicated parametrizations.
- Even in "simple" factor analysis models, there is still lots to explore ....
- Papers:
  - Sturma, Kranzlmueller, Portakal, Drton (2025). Matching Criterion for Identifiability in Sparse Factor Analysis. arXiv preprint arXiv:2502.02986.
  - Drton, Grosdos, Portakal, Sturma (2023). Algebraic Sparse Factor Analysis. arXiv preprint arXiv:2312.14762.





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